Graphical CSS Code Transformation Using ZX Calculus

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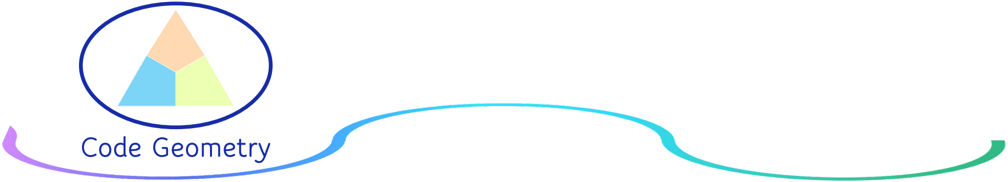
[6]Perimeter Institute for Theoretical Physics

IQC Student Seminar 2023

Can we unify the representations of a CSS code in one graphical language?



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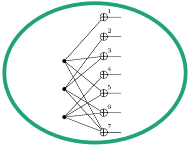
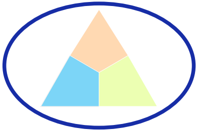


Tanner Graphs



Code Geometry

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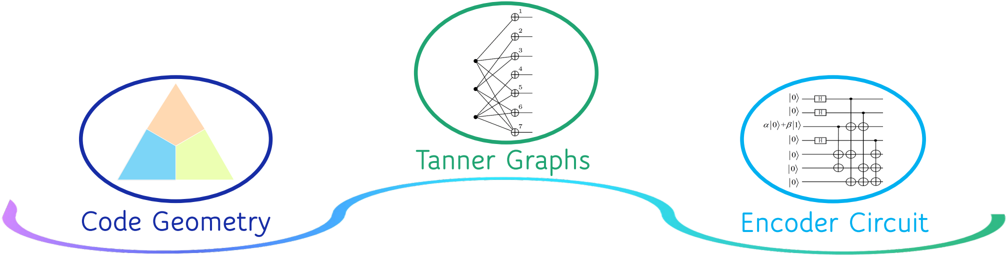

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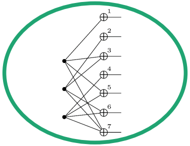

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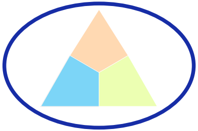

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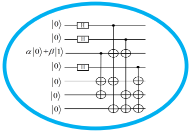

Encoder Circuit

YES

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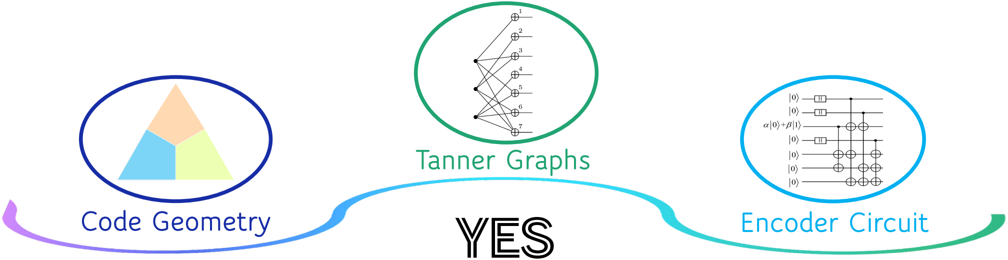
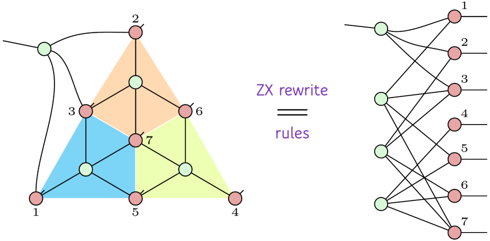


The ZX-calculus is an intuitive graphical language for quantum computation[1].

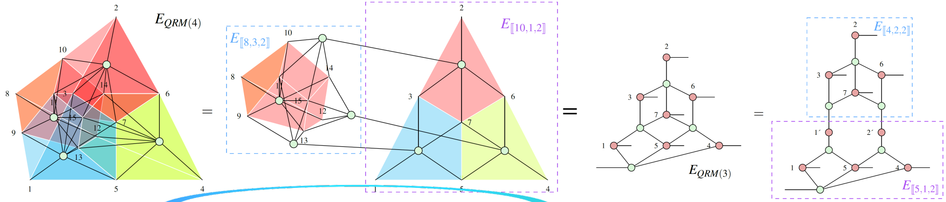


Any CSS encoder has a phase-free ZX normal form which corresponds to both the Tanner graph and geometry[2].

[1] Coecke, B., & Kissinger, A. (2018). Picturing quantum processes: A first course on quantum theory and diagrammatic reasoning. In *Diagrammatic Representation and Inference: 10th International Conference, Diagrams 2018, Edinburgh, UK, June 18-22, 2018, Proceedings 10* (pp. 28-31). Springer International Publishing.
[2] Kissinger, A. (2022). Phase-free ZX diagrams are CSS codes (... or how to graphically grok the surface code). In 19th International Conference on Quantum Physics and Logic.

Graphical transformation between CSS codes

Code Morphing



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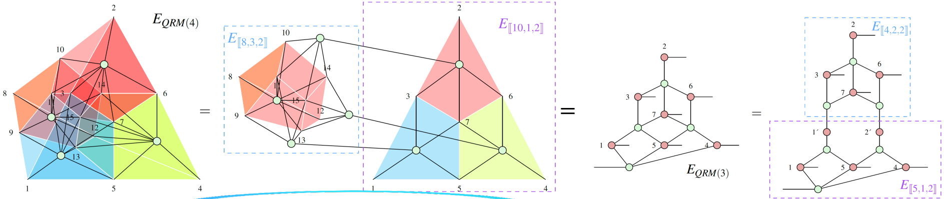

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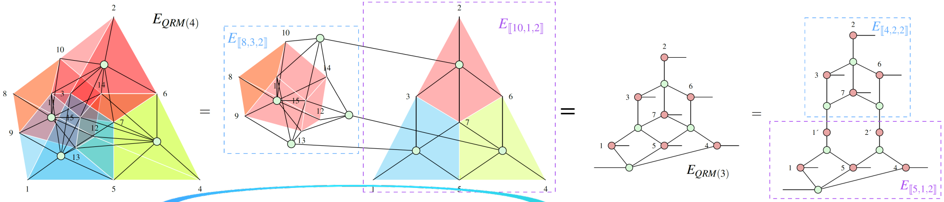

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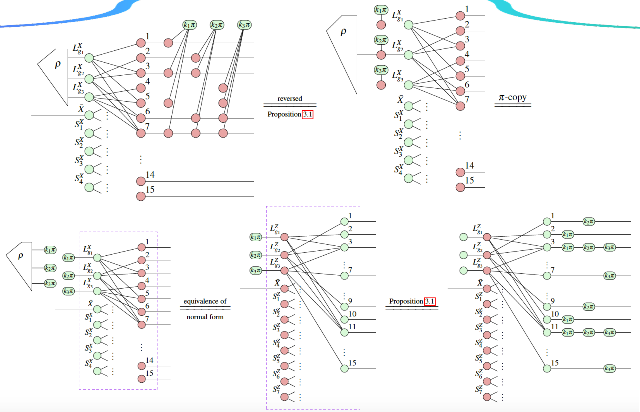

Subsystem Code Gauge Fixing

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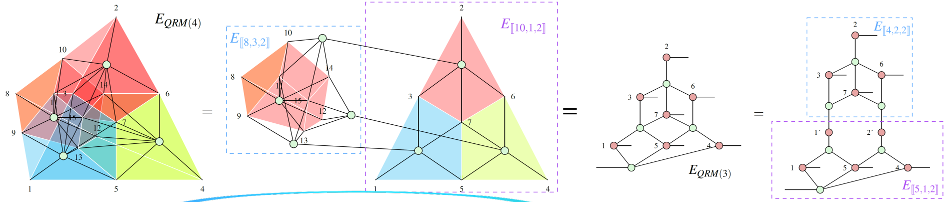
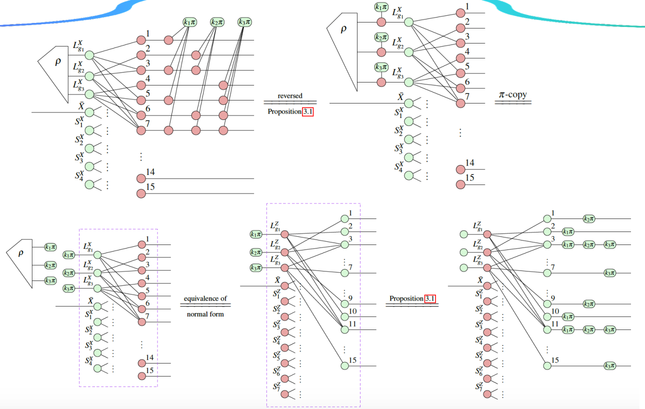


Code Switching

[4] Anderson, J. T., Duclos-Cianci, G., & Poulin, D. (2014). Fault-tolerant conversion between the steane and reed-muller quantum codes. Physical review letters, 113(8), 080501.

[5] Quan, D. X., Zhu, L. L., Pei, C. X., & Sanders, B. C. (2018). Fault-tolerant conversion between adjacent Reed–Muller quantum codes based on gauge fixing. Journal of Physics A: Mathematical and Theoretical, 51(11), 115305.



Subsystem Code Gauge Fixing

[6] Paetznick, A., & Reichardt, B. W. (2013). Universal fault-tolerant quantum computation with only transversal gates and error correction. Physical review letters, 111(9), 090505.

[7] Vuillot, C., Lao, L., Criger, B., Almud´ever, C. G., Bertels, K., & Terhal, B. M. (2019). Code deformation and lattice surgery are gauge fixing. New Journal of Physics, 21(3), 033028.

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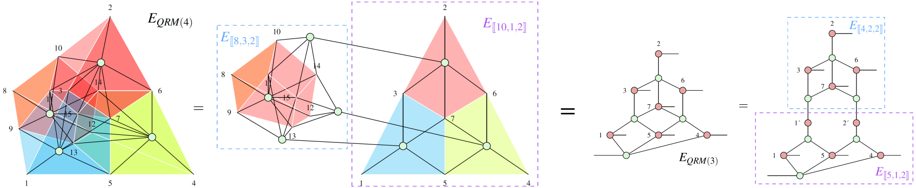

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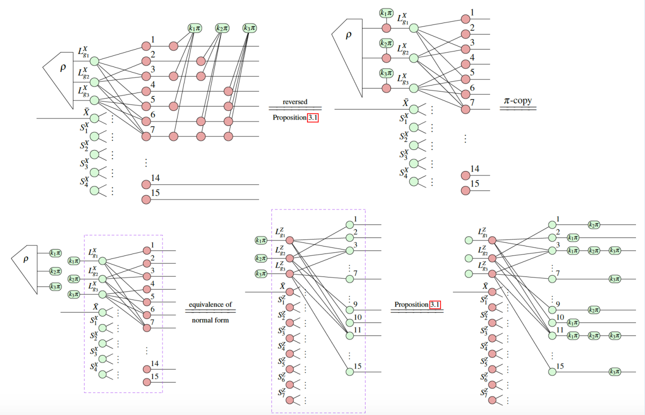
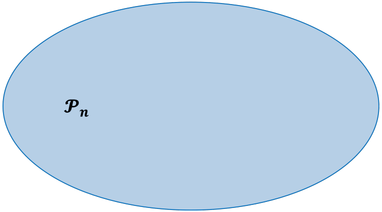


Subsystem Code Gauge Fixing

Stabilizer Code

Consider three groups of Pauli operators.

1. Pauli group on n qubits: $\mathcal{P}_n = \{i^c \left(\bigotimes_{i=1}^{n} P_i\right);\ P_i \in \{X, Y, Z, I\},\ 0 \leq c \leq 3\}$.



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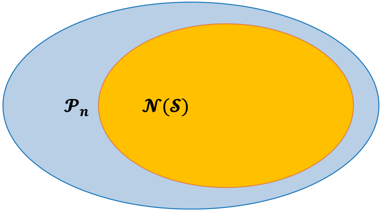
2. Stabilizer group: $\mathcal{S} = \langle M_1, M_2, \ldots, M_{n-k}\rangle$, $-I \notin \mathcal{S}$. $\mathcal{S} \subset \mathcal{P}_n$. $\mathcal{S}$ Abelian.



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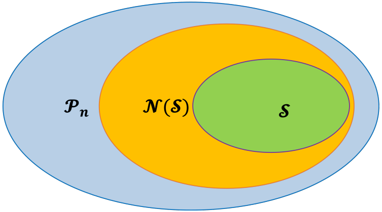

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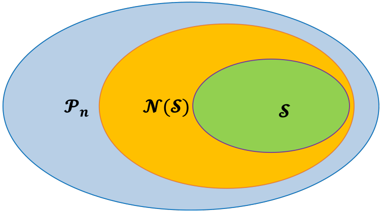

Definition

Stabilizer codes are a class of quantum error-correcting codes. Its code space $\mathcal{C}$ is the joint $+1$ eigenspace of $\mathcal{S}$.

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Code Space

$\lvert\overline{\psi}\rangle$ is called a *codeword* in $\mathcal{C}$, where

$$\mathcal{C} := \{n\text{-qubit state}\,\lvert\overline{\psi}\rangle\,;\, M\lvert\overline{\psi}\rangle = \lvert\overline{\psi}\rangle\,, \forall M \in \mathcal{S}\}$$

There are three important parameters for a stabilizer code: $[[n, k, d]]$.

- n is the number of physical qubits.
- k is the number of logical (or encoded) qubits.
- d is the code distance.

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Example

Consider $\mathcal{S} = \langle XX, ZZ\rangle$ on two qubits. Then $\mathcal{C} = \left\{\frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)\right\}$.

Logical Operators

Consider the centralizer of $\mathcal{S}$,

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- $\overline{X_1}, \overline{Z_1}, \ldots, \overline{X_k}, \overline{Z_k} \in \mathcal{N}(\mathcal{S})/\mathcal{S}$, up to the generators of $\mathcal{S}$.

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  They are anticommuting Pauli pairs acting non-trivially on $|\overline{\psi}\rangle$.
- All other operators in $\mathcal{P}_n$ anti-commute with at least one element in $\mathcal{S}$ and map a codeword $|\overline{\psi}\rangle$ onto a state outside the code space $\mathcal{C}$.

Fundamental Theorem of Stabilizer Theory

Theorem

If $\mathcal{S} \subset \mathcal{P}_n$ has m generators, then $\mathcal{C}$ is a 2^k dimensional subspace of $(\mathbb{C}^2)^{\otimes n}$, $k = n - m$.

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Example: Four-qubit code [[4, 2, 2]]

$$\mathcal{S} = \langle XXXX, ZZZZ \rangle$$

- What is the dimension of the code space?

Code Distance

Definition

Let d be the distance of a stabilizer code $\mathcal{C}(\mathcal{S})$, $|P|$ denotes the weight of $P \in \mathcal{P}_n$, the number of physical qubits on which P acts nontrivially. Then

$$d := \min_{P \in \mathcal{N}(\mathcal{S})/\mathcal{S}} |P|.$$

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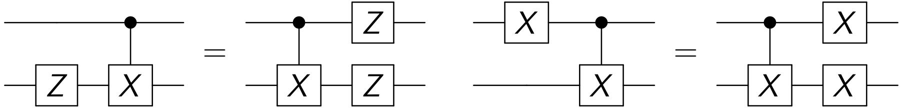
- Find pairs of mutually anti-commuting Paulis which commute with $XXXX$, $ZZZZ$.
- What is the code distance?

Fault-tolerant Technique: Transversality



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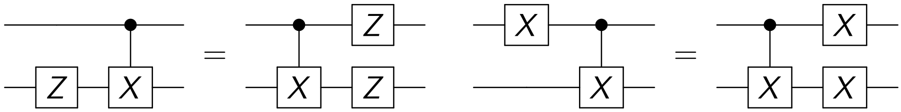

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A transversal logical operator is **NOT** implemented by any multi-qubit physical operation acting on the same code block.

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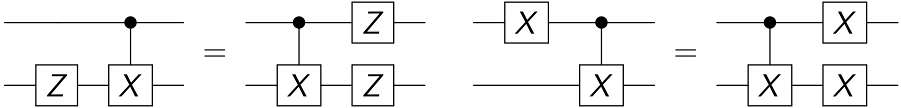

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- Transversality prevents any errors from spreading within a block, so a single physical error cannot cause a whole block of codes to go bad.

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Calderbank-Shor-Steane (CSS) Codes

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CSS codes are stabilizer codes whose stabilizer generators are defined by two orthogonal binary matrices G, H, $GH^T = 0$. Moreover,

$$\mathcal{S} = \langle G^X, H^Z \rangle.$$

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- $GH^T = 0$ implies that each X generator overlaps with a Z generator in an even number of places.
- Example: The $[\![7, 1, 3]\!]$ Steane code $\mathcal{S} = \langle S_1^X, S_2^X, S_3^X, S_1^Z, S_2^Z, S_3^Z \rangle$.

$$\begin{aligned} S_1^X &= X \otimes I \otimes X \otimes I \otimes X \otimes I \otimes X \\ &= X_1 X_3 X_5 X_7 \\ S_2^X &= X_2 X_3 X_6 X_7, \;\; S_3^X = X_4 X_5 X_6 X_7 \end{aligned}$$

$$\begin{aligned} S_1^Z &= Z \otimes I \otimes Z \otimes I \otimes Z \otimes I \otimes Z \\ &= Z_1 Z_3 Z_5 Z_7 \\ S_2^Z &= Z_2 Z_3 Z_6 Z_7, \;\; S_3^Z = Z_4 Z_5 Z_6 Z_7 \end{aligned}$$

Record Stabilizers for CSS Codes

Definition

Let M be a $k \times n$ binary matrix and $T \in \mathcal{X}, \mathcal{Z}$, then

$$M^T := \left\{ \bigotimes_{j=1}^{n} T^{[M]_{ij}} : 1 \leq i \leq k \right\} \subset \mathcal{P}^{\otimes n}.$$

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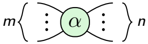

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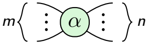

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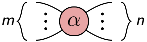

$$:= \; |+\rangle^{\otimes n} \langle +|^{\otimes m} \; + \; e^{i\alpha} \, |-\rangle^{\otimes n} \langle -|^{\otimes m} \, .$$

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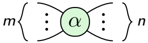

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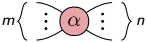
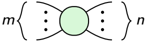
  - X spiders, which sum over the eigenbasis of the X operator:



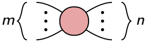
$$:= |+\rangle^{\otimes n} \langle +|^{\otimes m} + e^{i\alpha} |-\rangle^{\otimes n} \langle -|^{\otimes m}.$$

Definition

A ZX diagram is *phase-free* if its spiders have no phases.



$$:= |0\rangle^{\otimes n} \langle 0|^{\otimes m} + |1\rangle^{\otimes n} \langle 1|^{\otimes m}$$

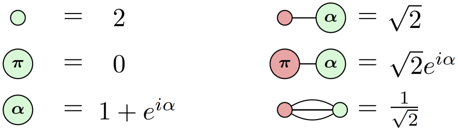


$$:= |+\rangle^{\otimes n} \langle +|^{\otimes m} + |-\rangle^{\otimes n} \langle -|^{\otimes m}$$

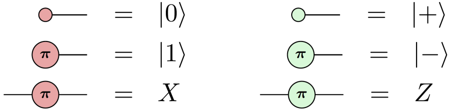
The ZX Calculus is Universal

Any linear map from *m* to *n* qubits corresponds exactly to a ZX diagram.

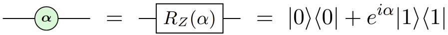
- A ZX diagram with 0 input and output represents a scalar.



- A ZX diagram with 0 input and 1 output represents a state.

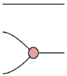


- A ZX diagram with the same number of inputs and outputs represents a unitary.

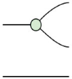


Represent the CNOT Gate in ZX

Proof. Horizontally composing the two diagrams below,

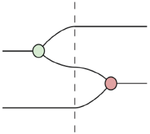


$$= \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \otimes \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{pmatrix} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 \end{pmatrix}$$

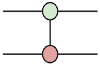


$$= \begin{pmatrix} 1 & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & 1 \end{pmatrix} \otimes \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

we get



$$= \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

Therefore, $CNOT = \sqrt{2}$ 

The ZX Calculus is Complete

If two ZX diagrams represent the same linear map, then there should be a sequence of rewrites that transforms one diagram into the other.

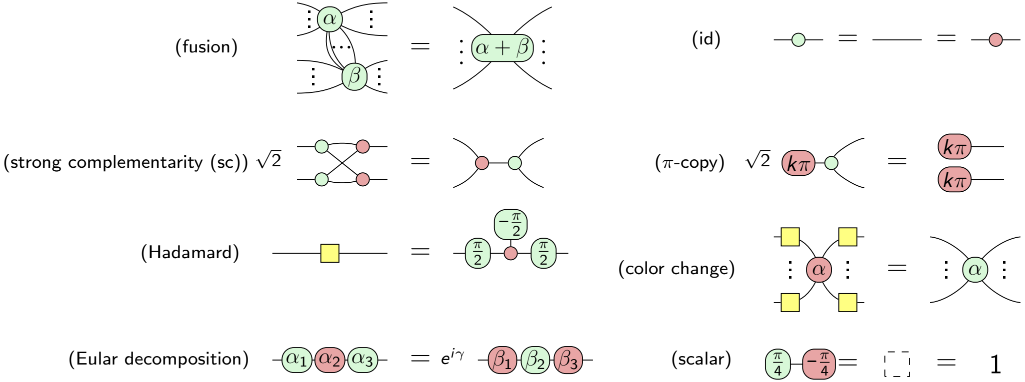
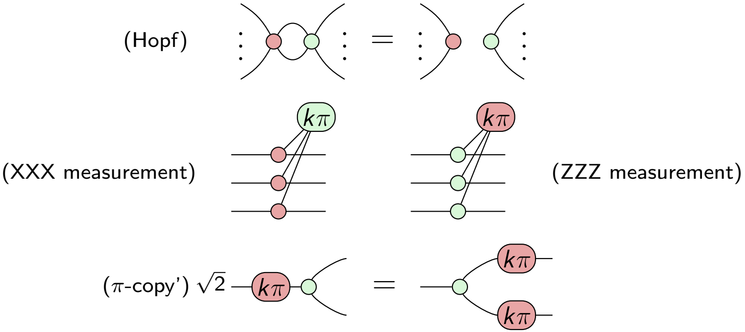


Figure: The minimal complete rule set for ZX calculus.

[3] Vilmart, R. (2019, June). A near-minimal axiomatisation of zx-calculus for pure qubit quantum mechanics. In 2019 34th Annual ACM/IEEE Symposium on Logic in Computer Science (LICS) (pp. 1-10). IEEE.

Additional ZX Rules

These rules are derivable from the minimal rule set. Used extensively in this work.



Phase-free ZX Diagrams are CSS Codes

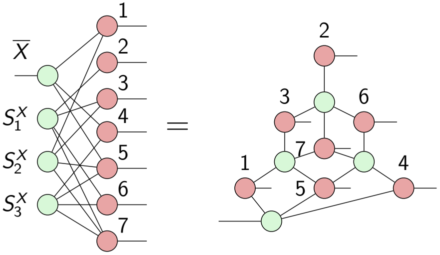
Consider an $[\![n, k, d]\!]$ CSS code with X-type stabilizers $\{S_1^X, S_2^X, \dots, S_{m'}^X\}$ and logical operators $\{\overline{X}_1, \overline{X}_2, \dots, \overline{X}_k\}$, it has a unique ZX normal form.

Example: The Steane code

- $n = 7,\ k = 1,\ d = 3.$
- 3 X-type stabilizers:

$$S_1^X = X_2X_3X_6X_7$$
$$S_2^X = X_1X_3X_5X_7$$
$$S_3^X = X_4X_5X_6X_7$$

- 1 logical X operator: $\overline{X} = X_1X_4X_5$



The Steane code encoder in ZX normal form.

[4]Kissinger, A. (2022). Phase-free ZX diagrams are CSS codes (... or how to graphically grok the surface code). arXiv preprint arXiv:2204.14038.

Subsystem Codes

Subsystem codes are stabilizer codes where some of the logical qubits are **NOT** used for information storage and processing. These logical qubits are called *gauge qubits*.

[5] David Kribs, Raymond Laflamme & David Poulin (2005): Unified and Generalized Approach to Quantum Error Correction. Physical Review Letters, 94.

Subsystem Codes

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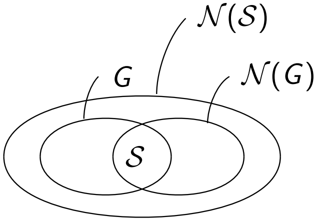
- G is an arbitrary subgroup of the Pauli group $\mathcal{P}$.
- $\mathcal{S} = \mathcal{N}(G) \cap G$, where $\mathcal{N}(G) = \{P \in \mathcal{P} : PM = MP, \forall M \in G\}$.
- $L_g = G \setminus \mathcal{S}$.



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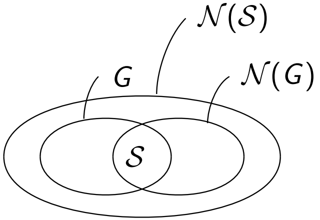


Definition

A subsystem code defined by G has a group $\mathcal{S}$ of stabilizers and a set L_g of gauge operators.

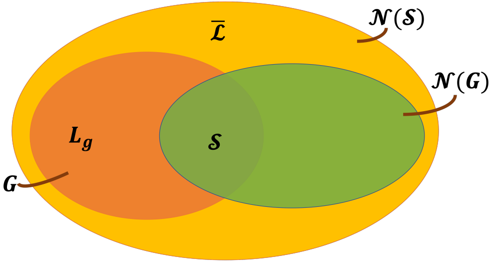
[5] David Kribs, Raymond Laflamme & David Poulin (2005): Unified and Generalized Approach to Quantum Error Correction. Physical Review Letters, 94.

CSS Subsystem Codes

Definition

CSS subsystem codes are subsystem codes whose stabilizers and gauge operators are either X-type or Z-type.

- When $L_g \neq \emptyset$, it contains pairs of anticommuting Pauli operators.



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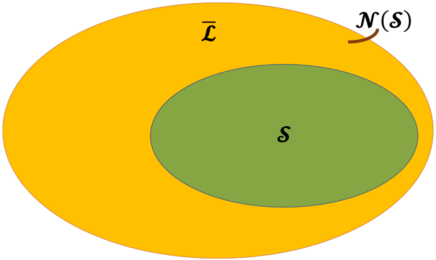
- When $L_g \neq \emptyset$, it contains pairs of anticommuting Pauli operators.
- When G is Abelian, $G = \mathcal{S}$ and $L_g = \emptyset$.



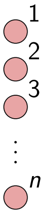
Normal Form for CSS Subsystem Codes

Example: For a CSS subsystem code $[[n, 1, d]]$, let $\overline{X}$ be the logical operator,

$$\mathcal{S}^X = \langle S_1^X, S_2^X, S_3^X \rangle, \quad L_g^X = \{L_{g_1}^X, L_{g_2}^X\}.$$

$L_{g_i}^X$ is a gauge operator acting non-trivially on the logical gauge qubit i. It acts trivially on other logical gauge qubits.

Step 1: For each physical qubit, introduce an X spider.



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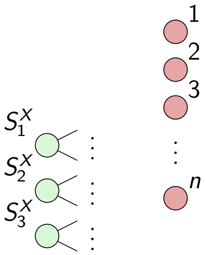
Step 2.1: For each X-type stabilizer S_i^X, introduce an Z spider. Connect it to all X spiders where S_i^X has support.



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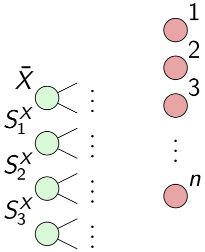
Step 2.2: For each X-type logical operator $\overline{X_j}$, introduce an Z spider. Connect it to all X spiders where $\overline{X_j}$ has support.

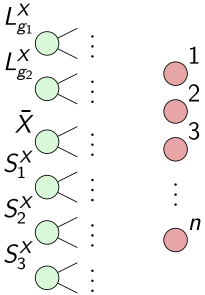


Normal Form for CSS Subsystem Codes

Example: For a CSS subsystem code $[[n, 1, d]]$, let $\overline{X}$ be the logical operator,

$$\mathcal{S}^X = \langle S_1^X, S_2^X, S_3^X \rangle, \quad L_g^X = \{L_{g_1}^X, L_{g_2}^X\}.$$

Step 2.3: For each X-type gauge operator $L_{g_t}^X$, introduce an Z spider. Connect it to all X spiders where $L_{g_t}^X$ has support.

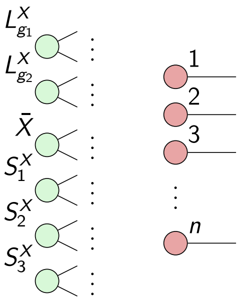


Normal Form for CSS Subsystem Codes

Example: For a CSS subsystem code $[[n, 1, d]]$, let $\overline{X}$ be the logical operator,

$$\mathcal{S}^X = \langle S_1^X, S_2^X, S_3^X \rangle, \quad L_g^X = \{L_{g_1}^X, L_{g_2}^X\}.$$

Step 3: Give each X spider an output wire.



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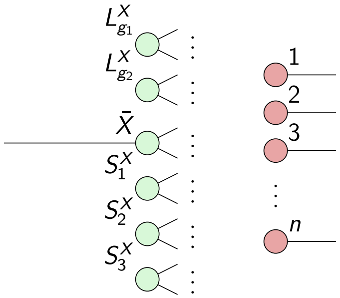
Step 4: For each Z spider representing $\overline{X_j}$, give it an output wire.



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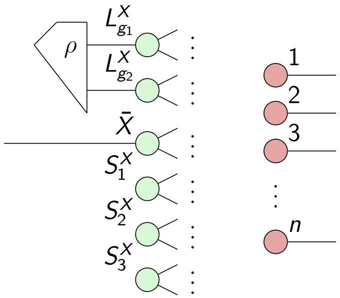
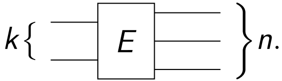
Step 5: For all Z spiders representing $L_{g_t}^X$, attach them to a joint arbitrary input state (i.e., a density operator ρ).



Pushing through the Encoder

For any $[\![n, k, d]\!]$ CSS code, its encoder map E is of the form:

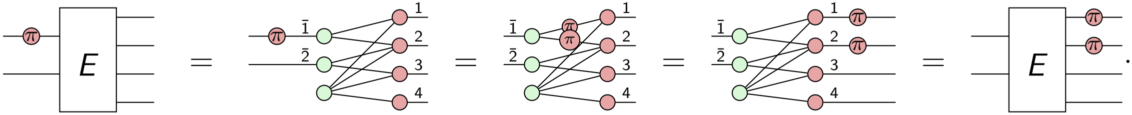


E is an isometry. $E^\dagger E = I$.

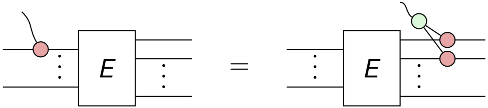
Lemma

In any CSS code, all $\overline{X_i}$ and $\overline{Z_i}$ must be multi-qubit Pauli operators.

Example: For the $[\![4, 2, 2]\!]$ code, $\overline{X_1} = X_1 X_2$:



Physically Implement a Logical Operator



Proposition

For any ZX diagram L on the left-hand side of E, one can write down a corresponding ZX diagram P on the right-hand side of E, such that $EL = PE$.

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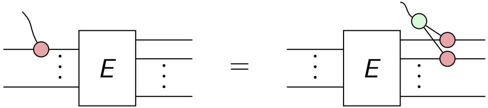

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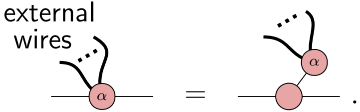
For any ZX diagram L on the left-hand side of E, one can write down a corresponding ZX diagram P on the right-hand side of E, such that $EL = PE$.

- Unfuse all spiders on logical qubit wires of L, whenever they are not phase-free or have more than one external wire.



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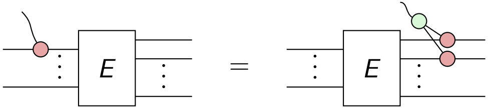

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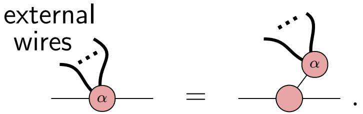
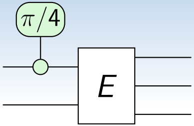
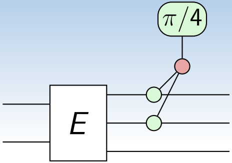


- For each X spider on logical qubit wires, rewriting E to be in ZX normal form and then applying the strong complementarity (sc) rule.



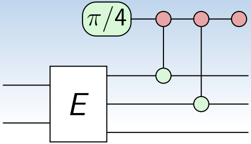
π/4
E

=

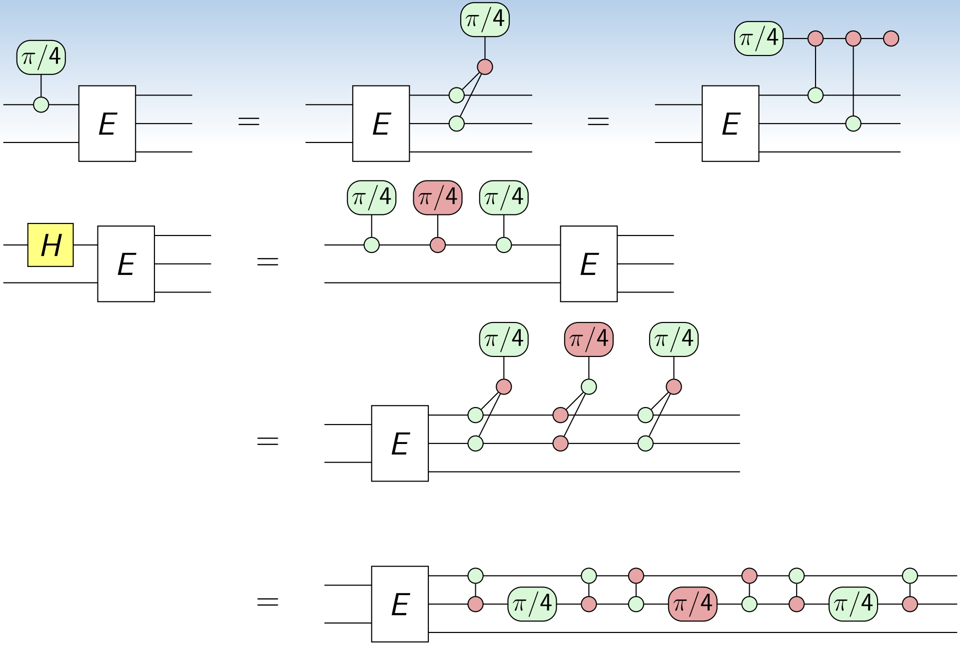


π/4
E

=

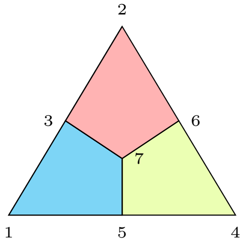


π/4
E

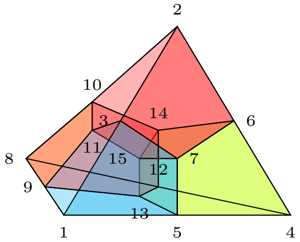

π/4
E
=
H

Switch between Two CSS Codes

	Steane Code: QRM(3)	Quantum Reed-Muller Code: QRM(4)
Code parameters	$[\![7, 1, 3]\!]$	$[\![15, 1, 3]\!]$
Logical operators	$\overline{X} = X_1 X_4 X_5$, $\overline{Z} = Z_1 Z_4 Z_5$	$\overline{X} = X_1 X_4 X_5$, $\overline{Z} = Z_1 Z_4 Z_5$
Transversal gates	$CX, S,$ **H**	$CX, S,$ **T**
Towards universality	Need transversal logical T	Need transversal logical H



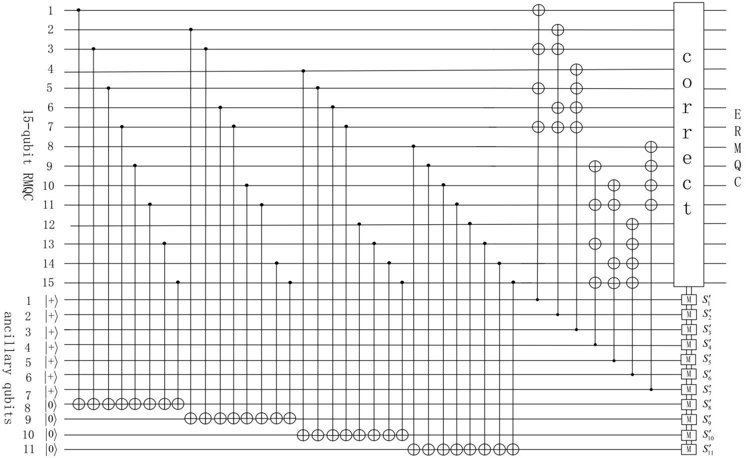
(a) QRM(3) as a 2D color code.



(b) QRM(4) as a 3D color code.

Code Switching[6,7]

- Codes with complementary fault-tolerant gate sets are switched between each other to realize a universal set of logical operations.
- Fault-tolerantly switch between QRM(3) and QRM(4)



[6] Anderson, J. T., Duclos-Cianci, G., & Poulin, D. (2014). Fault-tolerant conversion between the steane and reed-muller quantum codes. Physical review letters, 113(8), 080501.

[7] Quan, D. X., Zhu, L. L., Pei, C. X., & Sanders, B. C. (2018). Fault-tolerant conversion between adjacent Reed–Muller quantum codes based on gauge fixing. Journal of Physics A: Mathematical and Theoretical, 51(11), 115305.

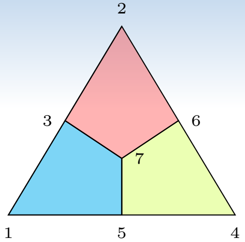
Steane Code & Quantum Reed-Muller Code

QRM(3) & QRM(4) are stabilizer codes defined by the stabilizers $\mathcal{S}_3$ & $\mathcal{S}_4$ respectively.

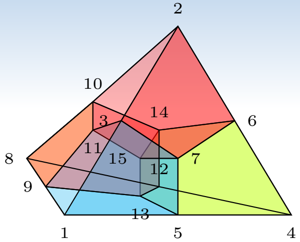
$$\mathcal{S}_3 = \langle M^X, M^Z \rangle, \qquad \mathcal{S}_4 = \langle N^X, N^Z, H^Z, T^Z \rangle, \qquad \text{where}$$

$$M = \begin{bmatrix} 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 \end{bmatrix}_{3\times 7}, \quad N = \begin{bmatrix} M & 0 & M \\ \mathbf{0} & 1 & \mathbf{1} \end{bmatrix}_{4\times 15}, \quad H = \begin{bmatrix} M & \mathbf{0} \end{bmatrix}_{3\times 15}$$

$$T = \begin{bmatrix} 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}_{3\times 15}$$



(a) QRM(3) as a 2D color code.



(b) QRM(4) as a 3D color code.

- $\mathcal{S}_3 = \langle M^X, M^Z \rangle$.
- Coloured face $\mapsto X/Z$ stabilizer generator.
- QRM(3) is self-dual.
- $\overline{X} = X_1X_4X_5$, $\overline{Z} = Z_1Z_4Z_5$.
- $d = 3$.

- $\mathcal{S}_4 = \langle N^X, N^Z, H^Z, T^Z \rangle$.
- Coloured face $\mapsto Z$ stabilizer generator.
- Coloured cell $\mapsto X/Z$ stabilizer generator.
- $\overline{X} = X_1X_4X_5$, $\overline{Z} = Z_1Z_4Z_5$.
- $d = 3$.

Subsystem Quantum Reed-Muller Code: SQRM

Definition

SQRM is defined by the gauge group:

$$G = \langle N^X, N^Z, H^Z, H^X, T^Z \rangle.$$

- The associated stabilizer group, gauge operators and logical operators are:

$$\mathcal{S}_S = \langle N^X, N^Z, H^Z \rangle_{11}, \qquad L_g = \langle H^X, T^Z \rangle_{6}, \qquad \overline{L} = \langle \overline{X}, \overline{Z} \rangle_{2}.$$

- For brevity, we will use $L_g^X = H^X$ and $L_g^Z = T^Z$.

$\Rightarrow$ SQRM has 1 logical qubit and 3 gauge qubits.

$\Rightarrow$ Alternatively, $\mathcal{S}_S$ stabilizes the $[\![15, 4, 3]\!]$ CSS code, with logical operators $\{L_g, \overline{L}\}$.

Extended Quantum Reed-Muller Code: EQRM

EQRM is defined by the stabilizer group $\mathcal{S}_E = \langle N^X, N^Z, H^Z, H^X \rangle$, where

$$M = \begin{bmatrix} 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 \end{bmatrix}_{3\times 7}, \qquad N = \begin{bmatrix} M & 0 & M \\ \mathbf{0} & 1 & \mathbf{1} \end{bmatrix}_{4\times 15}, \qquad H = \begin{bmatrix} M & \mathbf{0} \end{bmatrix}_{3\times 15}.$$

- Let $|\bar{0}\rangle$ and $|\bar{1}\rangle$ be the logical 0 and 1 encoded in QRM(3).
- Let $|\bar{\psi}\rangle = \alpha |\bar{0}\rangle + \beta |\bar{1}\rangle$ be the single-qubit logical information encoded in QRM(3).

Lemma

An EQRM codeword $|\bar{\Phi}\rangle$ can be decomposed as

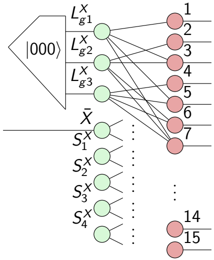
$$|\bar{\Phi}\rangle = |\bar{\psi}\rangle \otimes |\phi\rangle, \text{ where}$$

$|\phi\rangle = \frac{1}{\sqrt{2}}(|0\rangle |\bar{0}\rangle + |1\rangle |\bar{1}\rangle)$.

Gauge Fixing SQRM

Lemma

When the three gauge qubits are in the $|\overline{000}\rangle$ state, SQRM is fixed to QRM(4).



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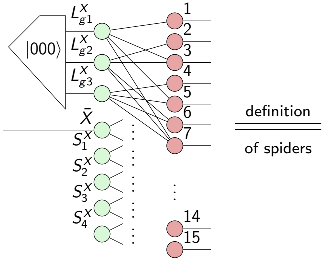

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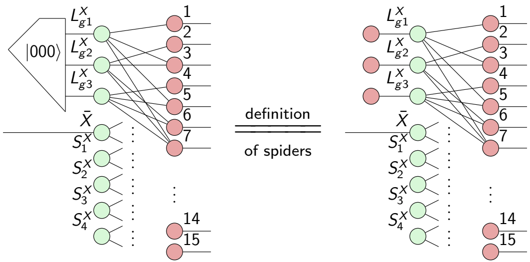

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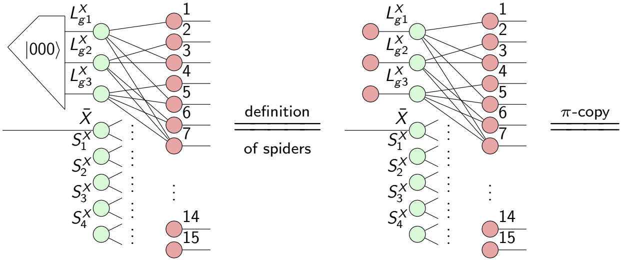

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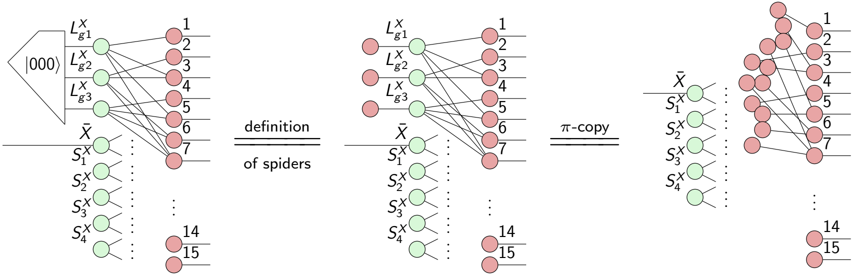

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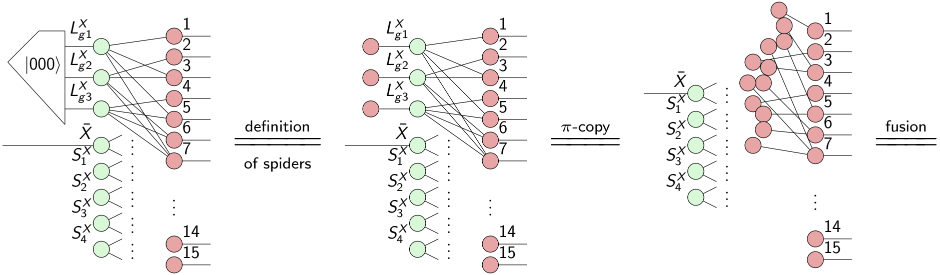

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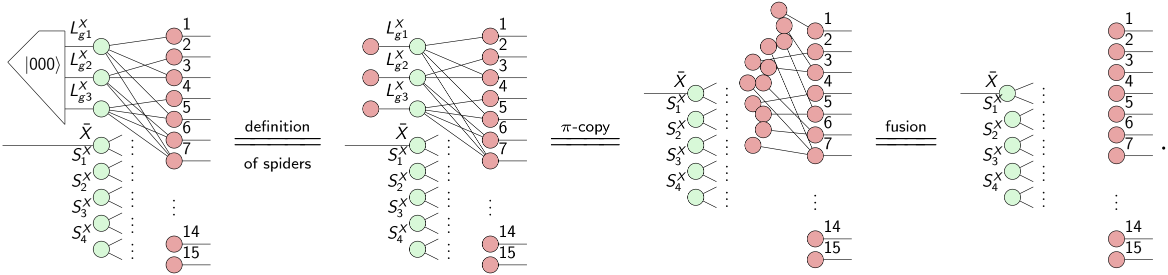

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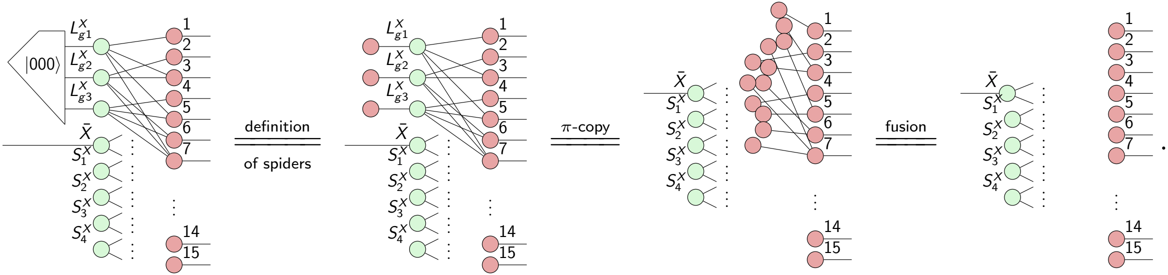

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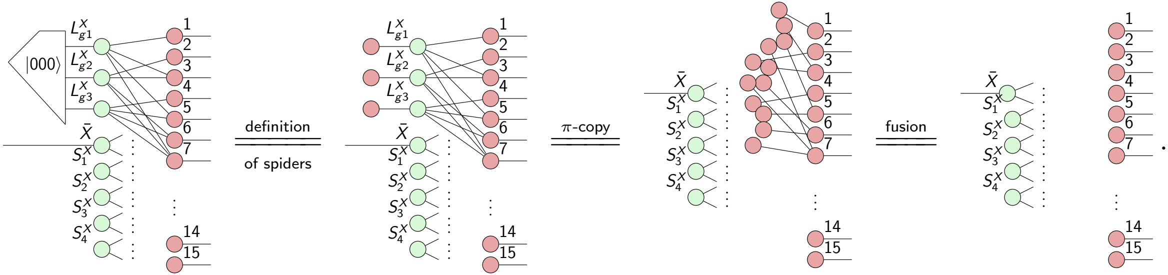

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$\Rightarrow$ Start with the XZ normal form of the SQRM encoder.

Forward & Backward Switching

- Forward: EQRM $\Rightarrow$ QRM(4) $\mathcal{S}_4$

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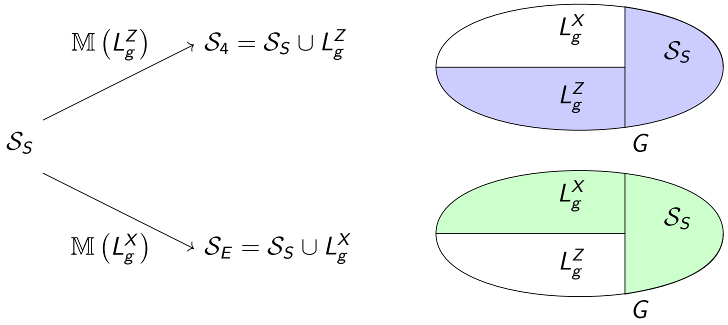

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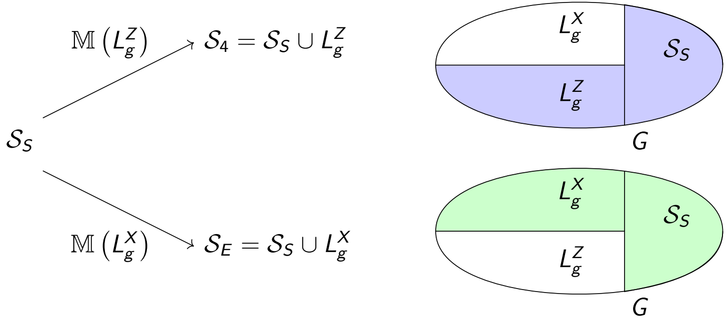

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Gauge Fixing SQRM Realizes EQRM

Step 1: : Measure a commuting subset of gauge operators. E.g., measure three X-type gauge operators $L^X_{g_i}$ and obtain the corresponding outcomes $k_1, k_2, k_3 \in \{0, 1\}$:

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$$Z \left|0\right\rangle = \left|0\right\rangle , \; Z \left|1\right\rangle = -\left|1\right\rangle$$

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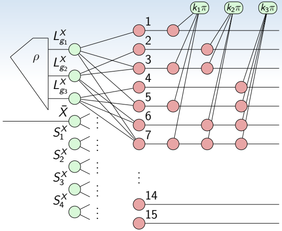

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$k_3\pi$
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$L^X_{g_2}$
$L^X_{g_3}$
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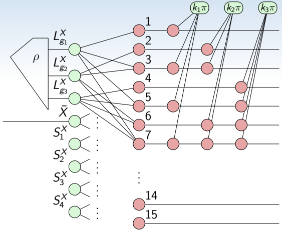

reversed

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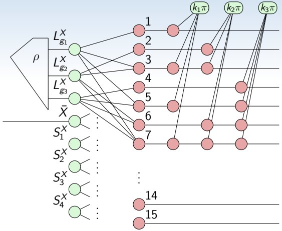

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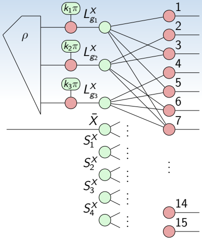

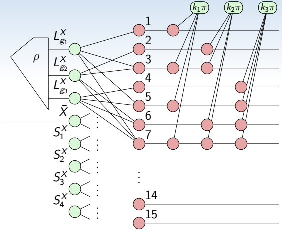

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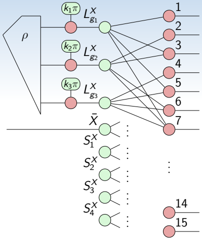

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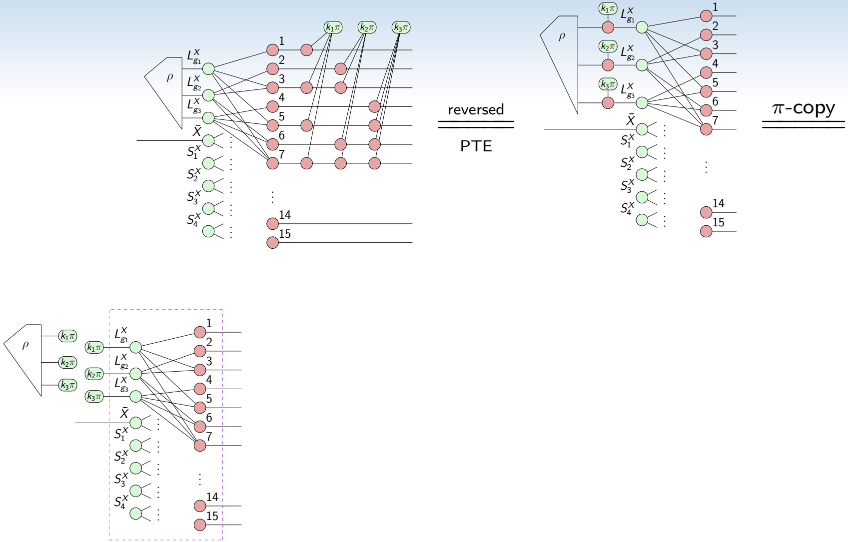

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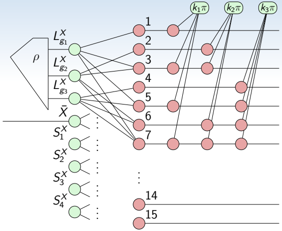

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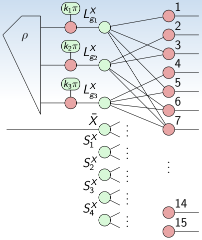

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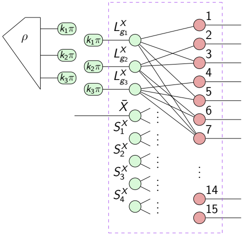

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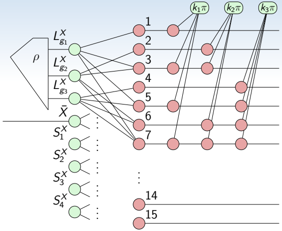

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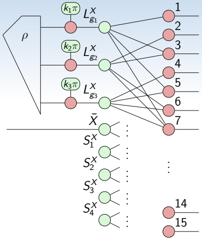

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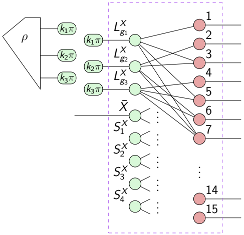

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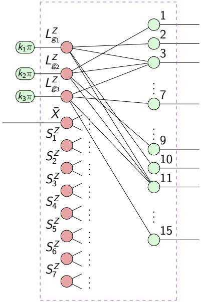

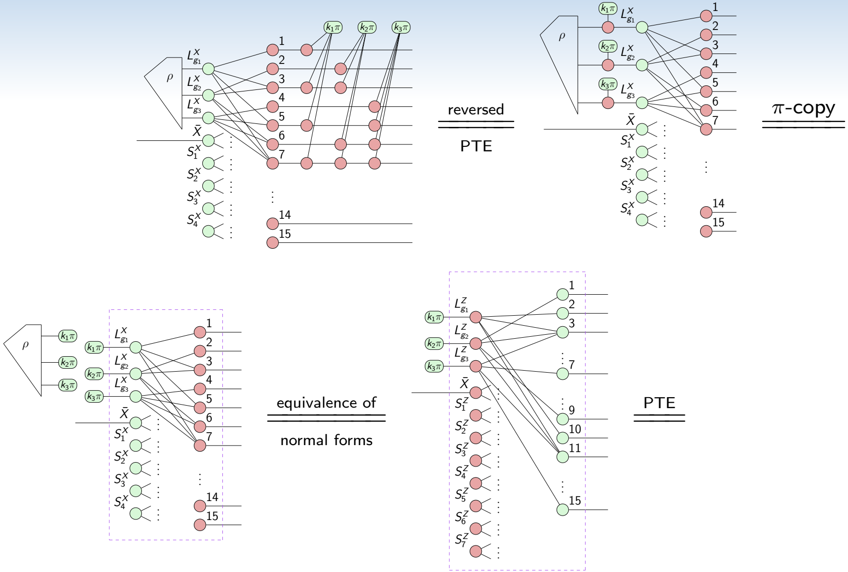
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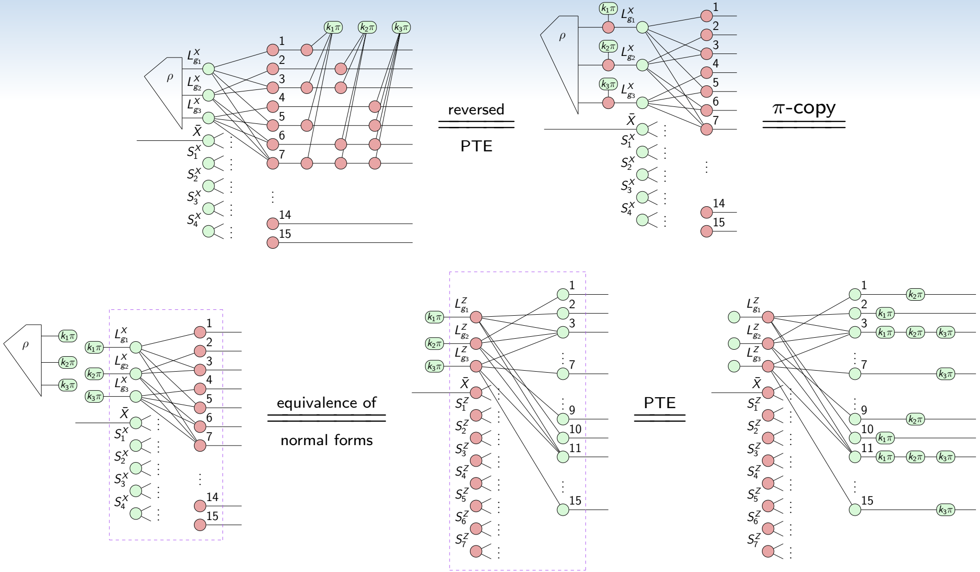
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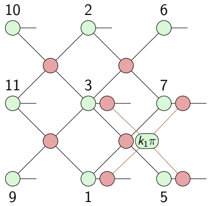

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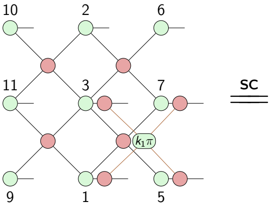

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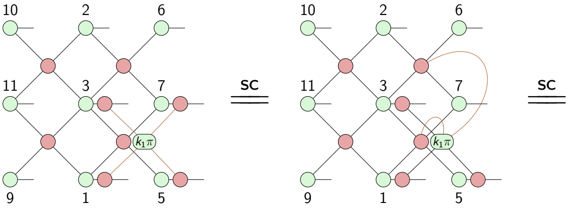

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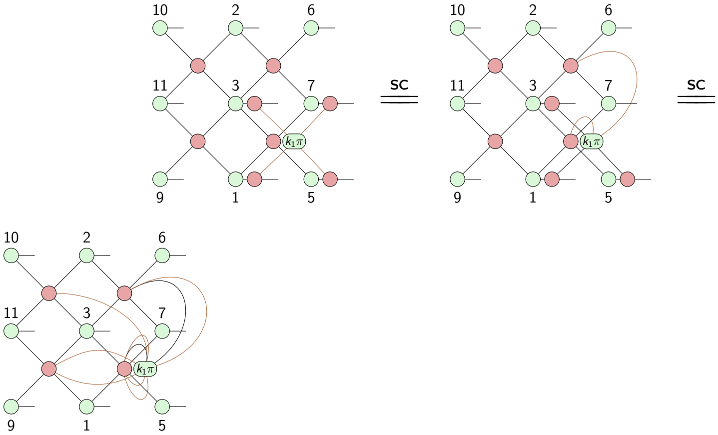

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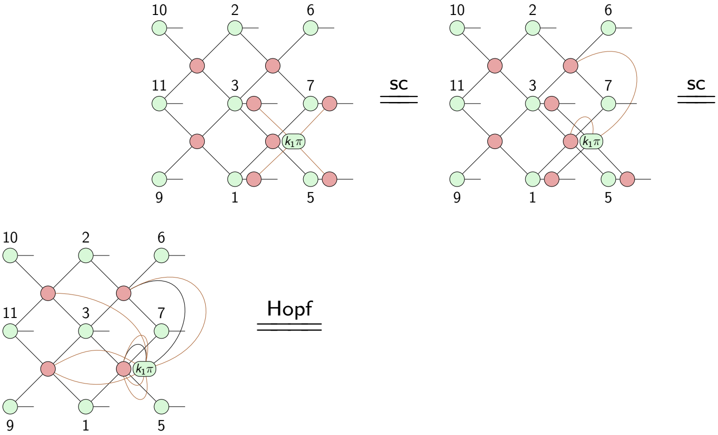

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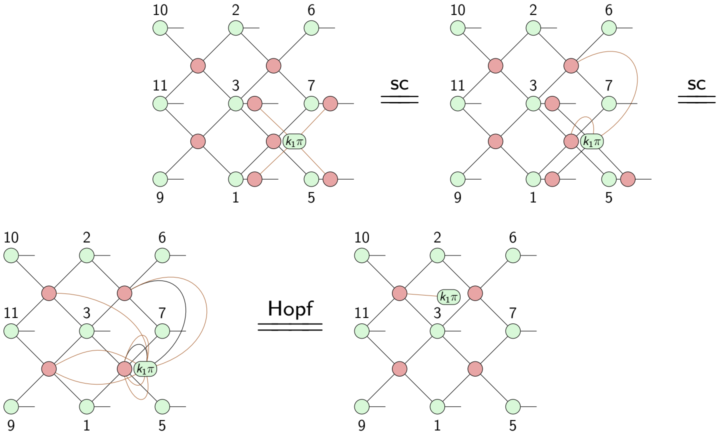

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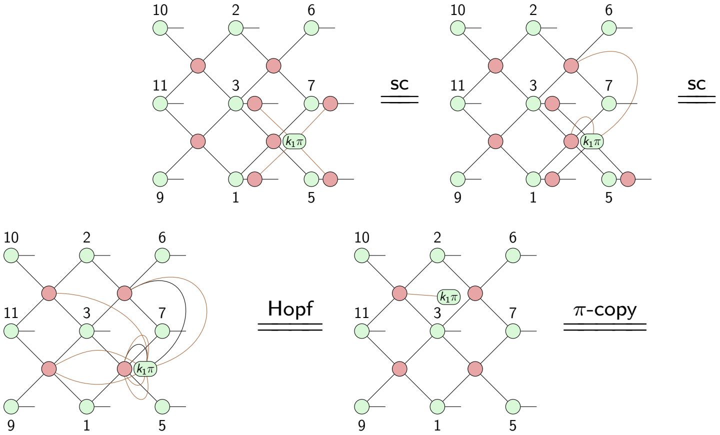

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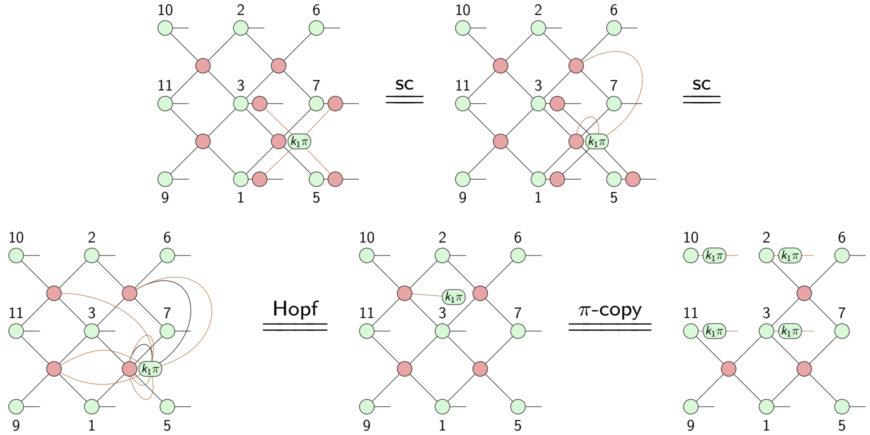

Construct SQRM from QRM(4) & EQRM

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Open Problems

Through the lens of ZX calculus, we will

- **Present** CSS code deformations.

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[13] Khesin, A. B., Lu, J. Z., & Shor, P. W. (2023). Graphical quantum Clifford-encoder compilers from the ZX calculus. arXiv preprint arXiv:2301.02356.

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[12] Bombin, H., Litinski, D., Nickerson, N., Pastawski, F., & Roberts, S. (2023). Unifying flavors of fault tolerance with the ZX calculus. arXiv preprint arXiv:2303.08829.

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Open Problems

Through the lens of ZX calculus, we will

- **Present** CSS code deformations.
- **Understand** code concatenation.
- **Derive** new good QECCs from the existing QECCs.
- **Unify** fault-tolerant protocols for stabilizer codes.

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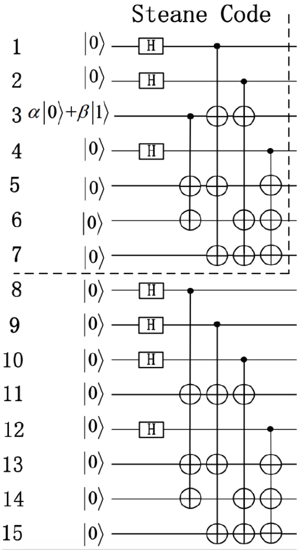
[13] Khesin, A. B., Lu, J. Z., & Shor, P. W. (2023). Graphical quantum Clifford-encoder compilers from the ZX calculus. arXiv preprint arXiv:2301.02356.

Thank you!



$|\bar{\psi}\rangle = \alpha|\bar{0}\rangle_3 + \beta|\bar{1}\rangle_3$

$\mathbb{S}_3$ stabilizes $|\bar{\psi}\rangle$

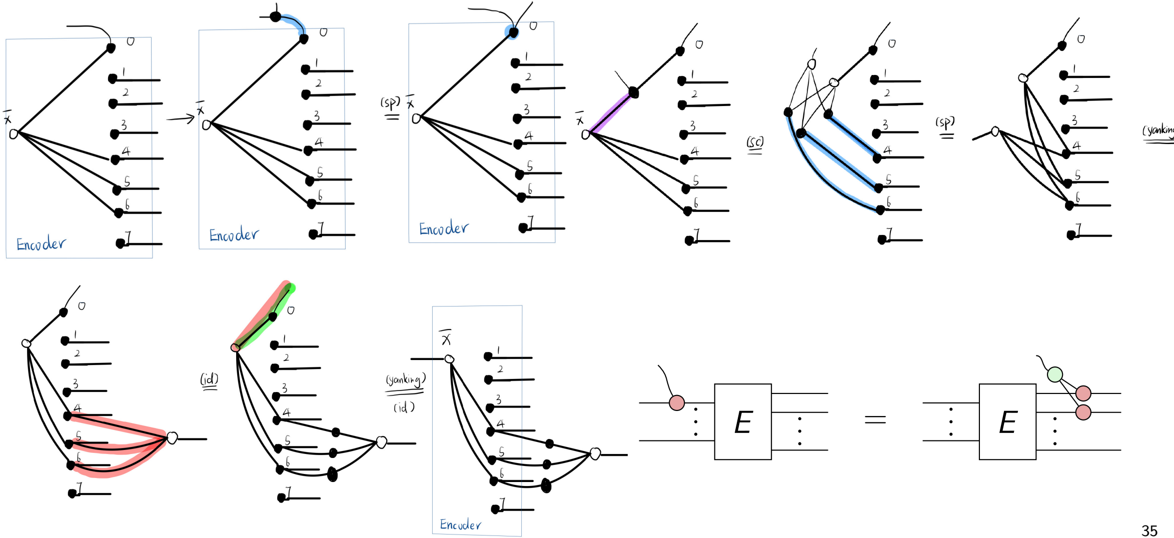
$|\phi\rangle = \frac{1}{\sqrt{2}}(|0\rangle|\bar{0}\rangle_3 + |1\rangle|\bar{1}\rangle_3)$

No information is stored on the last eight qubits.

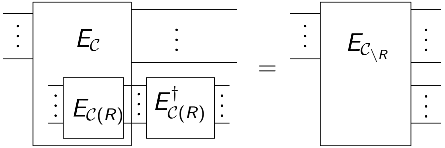
$|\bar{\Phi}\rangle = |\bar{\psi}\rangle|\phi\rangle$

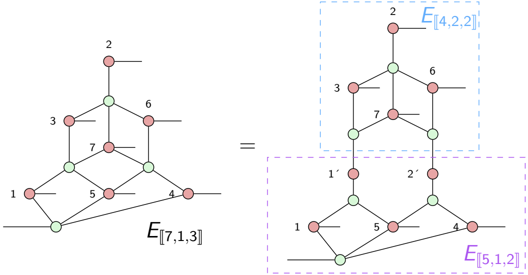
$\mathbb{S}_E$ stabilizes $|\bar{\Phi}\rangle$

Example: $\overline{X} = X_4X_5X_6$



Code Morphing[10]





[10] Michael Vasmer & Aleksander Kubica (2022): Morphing Quantum Codes. PRX Quantum 3(3), doi:10.1103/prxquantum.3.030319